# ASSIGNMENT 5 - FUNCTIONAL PROGRAMMING 

COMP 3010 - ORGANIZATION OF PROGRAMMING LANGUAGES

## 1. Lisp \& Scheme

Exercise 1. Consider the Scheme definition:
(define (mystery x y)
(lambda (z) (x (y (x z)))))
(1) Translate the mystery function to a $\lambda$-calculus expression. HINT: You will need $\lambda_{\mathrm{s}}$ introducing the parameters $x$ and $y$ around the body of mystery.
(2) What happens when you evaluate
(mystery (lambda (x) (+ 1 x)) (lambda (y) (* 2 y)))
(3) What happens when you evaluate
((mystery (lambda (x) (+ 1 x)) (lambda (y) (* 2 y))) 5)
(4) What happens when you evaluate
(( (mystery (lambda (x) (+ 1 x)) (lambda (y) (* 2 y))) 5) 6)
Exercise 2. Remember the map function, which changes every element of a list using a given operation, is written in Scheme as

```
(define (map change-car xs)
    (cond [(null? xs) xs]
            [(cons? xs) (cons (change-car (car xs))
                            (map change-car (cdr xs)))]))
```

so that a list (list x y $\ldots \mathrm{z}$ ) is transformed like so

```
(map f (list x y ... z)) = (list (f x) (f y) ... (f z))
```

reduce, which compresses a list by replacing every cons with a chosen binary operation and the final null with a chosen constant, is written in Scheme as

```
(define (reduce change-cons change-null xs)
    (cond [(null? xs) change-null]
            [(cons? xs) (change-cons (car xs)
                            (reduce change-cons change-null (cdr xs)))]))
so that a list (cons \(x\) (cons y (cons... (cons z null)))) is transformed as
    (reduce fe (cons x (cons y (cons ... (cons z null))))
    \(=(f x\) (f y (f ... (f z e))))
    (1) What is the result of evaluating
        (map (lambda (x) (* x x)) (list \(\left.12 \begin{array}{llll} & 3 & 4\end{array}\right)\) )
    (2) What is the result of evaluating
        (reduce + O (list 1234 5))
    (3) What is the result of evaluating
        (reduce + 0 (map (lambda (x) (* x x)) (list 12345 )) )
```

[^0](4) (Multiple Choice) Consider this definition of the function $f$ :
(define (f xs)

```
(reduce + O (map (lambda (x) (* x x)) xs)))
```

Which of the following alternate definitions of $f$ is equivalent to the one above that used map and reduce?
(a) (define (f xs)
(cond [(null? xs) 0]
[(cons? xs) (* (+ (car xs) (car xs)) (f (cdr xs)))]))
(b) (define (f xs)
(cond [(null? xs) 0]
[(cons? xs) (+ (* (car xs) (car xs)) (f (cdr xs)))]))
(c) (define (f xs)
(cond [(null? xs) 0]
[(cons? xs) (+ (car xs) (f (cdr xs)))]))
(d) (define (f xs)
(cond [(null? xs) 0]
[(cons? xs) (* (car xs) (f (cdr xs)))]))

## 2. ML Family

Exercise 3. Do Concepts In Programming Languages Exercise 5.3 on Nonlinear Pattern Matching (page 123).

Note, for parts (a) and (b), you can write the described functions in SML syntax as asked by the exercise, $O R$ in your choice of Ruby, Python, or C syntax.

Exercise 4. Do Concepts In Programming Languages Exercise 5.7 on Disjoint Unions (page 125).

Exercise 5. In SML, all references must point to real values in the heap. In other words, SML does not support implicit null pointers in place of a reference. Instead, the SML data type declaration

```
datatype 'a option = NONE | SOME of 'a;
```

defines the generic type 'a option of references which could either point to nothing (represented by the NONE constructor containing no data) or point to some actual ' a in the heap (represented by the SOME constructor containing a value of type 'a).

For example, the integer division operation x div y will raise an exception when the divisor $y$ is 0 . A safe version of division, which never raises an exception, can be written in SML as

```
fun safe_div(x, 0) = NONE
    | safe_div(x, y) = SOME (x div y);
```

which takes a pair of ints and returns an int option.
(1) What is the difference between the result of evaluating 10 div 0 versus safe_div(10, 0)?
(2) What is the difference between the result of evaluating 10 div 5 versus safe_div(10, 5)?
(3) What happens when you try to evaluate 2 * (10 div 5)? What happens when you try to evaluate $2 *$ (safe_div $(10,5))$ ?


[^0]:    Date: Spring 2023.

